

Lecture 18

Recall: $G \sim G(n, p)$ means that G is a graph on n (labeled) vertices, where each edge appears indep. w/ prob p . Erdős-Rényi
! & !. are diff. graphs

Monotone graph properties have sharp thresholds:

$$P\{G \sim G(n, p) \text{ has property } P\} \approx \begin{cases} 0 & \text{if } p \ll t(n) \\ 1 & \text{if } p \gg t(n) \end{cases}$$

Example: $P = \{G \text{ has } \geq 1 \text{ edge}\}$

↪ monotone property

Claim: $t(n) = \frac{1}{n^2}$

↳ why?

Let $X = \# \text{ edges in } G$, $P_n = \frac{c}{n^2}$

$$\begin{aligned} P\{G \text{ has } 0 \text{ edges}\} &= P\{X=0\} \\ &= (1 - P_n)^{\binom{n}{2}} \quad \text{why?} \\ &\rightarrow e^{-c/2} \text{ as } n \rightarrow \infty \\ &\approx \begin{cases} 1 & \text{for } c \ll 1 \\ 0 & \text{for } c \gg 1 \end{cases} \end{aligned}$$

Goal: For $P = \{G \text{ is connected}\}$, show $t(n) = \frac{\log n}{n}$

↳ Thm (Erdős-Rényi)

Let $P_n = \frac{\lambda \log n}{n}$, λ fixed

if $\lambda > 1$, then $P\{G \sim G(n, P_n) \text{ is connected}\} \rightarrow 1$ as $n \rightarrow \infty$

$\lambda < 1$, then $P\{G \sim G(n, P_n) \text{ is connected}\} \rightarrow 0$ as $n \rightarrow \infty$

Lemma: If X is a r.v. then

$$P(X=0) \leq \frac{\text{Var}(X)}{(\mathbb{E}[X])^2}$$

Pf: $\text{Var}(X) = \mathbb{E}[X - \mathbb{E}[X]]^2 = P(X=0) \underbrace{\mathbb{E}[X - \mathbb{E}[X]]^2 \mid X=0}_{\mathbb{E}[X]^2} + P(X \neq 0) \underbrace{\mathbb{E}[X - \mathbb{E}[X]]^2 \mid X \neq 0}_{\geq 0}$

$$\geq P(X=0) (\mathbb{E}[X])^2$$

Case 1: $\lambda < 1$:

Will show: $P\{G \text{ has an isolated vertex}\} \rightarrow 1$

Note: $\{G \text{ has an isolated vertex}\} \subset \{G \text{ is disconnected}\}$ subset of

$$X = \sum_{i=1}^n I_i = \# \text{ isolated vertices where } I_i = \begin{cases} 0 & \text{if } i \text{ not isolated} \\ 1 & \text{if } i \text{ isolated} \end{cases}$$

$$I_i \sim \text{Bern}(q)$$

$$q := (1 - p_n)^{n-1} \quad \begin{matrix} \text{(n-1) edges must not be} \\ \text{placed} \\ \text{prob of an edge not} \\ \text{being placed} \end{matrix}$$

$$\mathbb{E}[X] = n \mathbb{E}[I_i] = nq$$

$$\text{Var}(X) = \sum_i \text{Var} I_i + \sum_{i \neq j} \text{Cov}(I_i, I_j)$$

$$= nq(1-q) + n(n-1) \text{Cov}(I_1, I_2)$$

due to the symmetric nature of the graphs

$$= \mathbb{E}[I_1, I_2] - \mathbb{E}[I_1]^2$$

$$= nq(1-q) + n(n-1) (\mathbb{E}[I_1, I_2] - \mathbb{E}[I_1]^2)$$

$P\{\text{vertices 1 \& 2 are isolated}\}$

$$= \frac{q^2}{1-p_n} - q^2$$

$$= nq(1-q) + n(n-1) \left(\frac{p_n q^2}{1-p_n} \right)$$

Why aren't we considering $\mathbb{E}[I_i]^2$



$$P\{G \text{ has no isolated vertices}\} = P\{X=0\} \leq \frac{nq(1-q) + n(n-1) \frac{p_n q^2}{1-p_n}}{n^2 q^2}$$

$$\leq \frac{(1-q)}{nq} + \frac{p_n}{1-p_n}$$

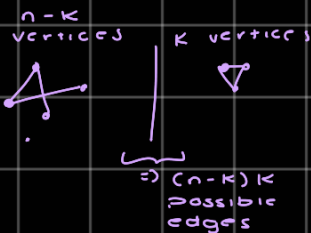
$\underbrace{nq}_{\rightarrow 0} \rightarrow 0$ bc $nq \rightarrow \infty$ $\underbrace{\frac{p_n}{1-p_n}}_{\rightarrow 0 \text{ as } n \rightarrow \infty}$

bc $\log(nq) \sim \log n^{(1-\lambda)}$

$\therefore P\{G \text{ has an isolated vertex}\} \rightarrow 1$

Case 2: $\lambda > 1$:

$$P\{G \text{ disconnected}\} = P \left\{ \bigcup_{k=1}^{n/2} \left\{ \text{set of } k \text{ vertices separated from the rest of } G \right\} \right\}$$



union bound

$$\leq \sum_{k=1}^{n/2} P\{\text{set of } k \text{ vertices separated from the rest of } G\}$$

$$\leq \sum_{k=1}^{n/2} \binom{n}{k} P\{\text{vertices } \{1, \dots, k\} \text{ can be separated from rest of } G\}$$

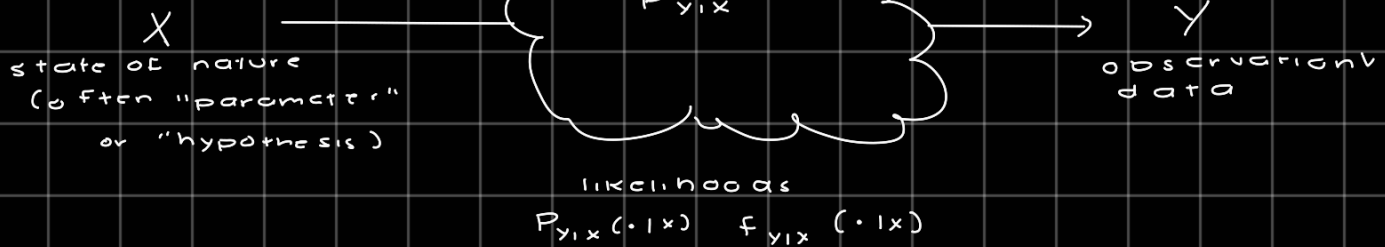
$$= \sum_{k=1}^{n/2} \binom{n}{k} (1-p_n)^{k(n-k)}$$

lots of tedious work $\rightarrow \square$

→ Now, onto module 3!

Inference: Want to infer X from knowing Y

• Basic Setup:



X may or may not be a r.v.

↳ If it is, then distribution of X is called a "prior"
(Bayesian inference)

↳ Example of inference when no prior is available

• Maximum Likelihood Estimation:

$$\text{MLE}(X | Y=y) := \underset{x}{\operatorname{argmax}} P_{y|x}(y|x) \quad (\text{break ties arbitrarily})$$

→ ex: model is $y \sim \mathcal{N}(x, \sigma^2)$

$$\text{MLE}(X | Y=y) = y$$